

#### Chapter 13 IPsec

### IPsec (IP Security)

- A collection of protocols used to create VPNs
- A network layer security protocol providing cryptographic security services that can support various combinations of authentication, integrity, access control, and confidentiality
- Allows creation of an encrypted tunnel between two private networks
- Supports authentication of the two ends of the tunnel
- Cannot directly encrypt non-IP traffic
- Can encrypt GRE tunnel containing non-IP data
- Comprises of IKE, ESP, and AH

#### Types of IPsec VPNs

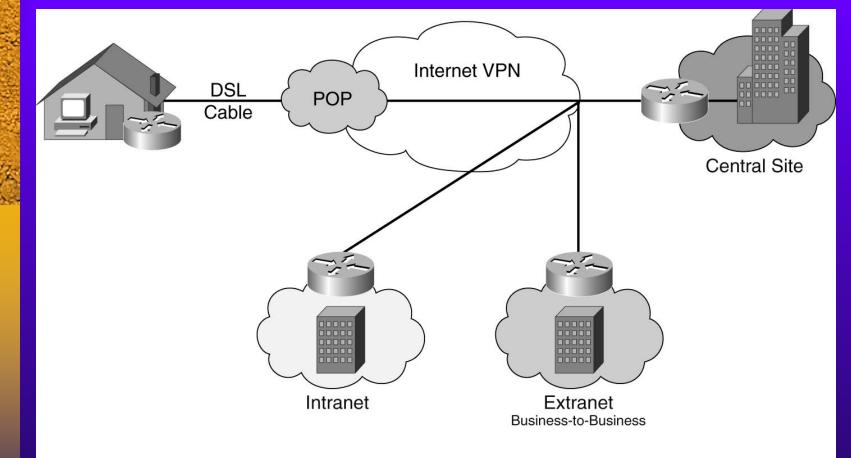
#### LAN-to-LAN or site-to-site

 Used to connect two private networks to form one combined virtual private network

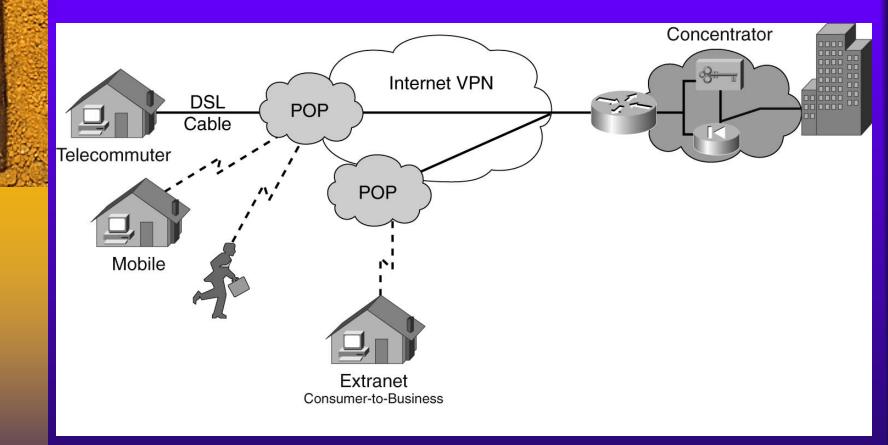
#### Remote-access client IPsec

 Used to allow road warriors to be part of the trusted network

### LAN-to-LAN and Site-to-Site IPsec



#### Remote-Access IPsec



#### **IPsec Protocol Suite**

#### Internet Key Exchange (IKE) protocol

- For negotiating security parameters and establishing authenticated keys
- Uses UDP port 500 for ISAKMP

Encapsulating Security Payload (ESP) protocol

- For encrypting, authenticating, and securing data
- IP protocol 50
- Authentication Header (AH) protocol
  - For authenticating and securing data
  - IP protocol 51

# IKE's Responsibilities in IPsec Protocol

- Negotiates IPsec tunnel characteristics between two IPsec peers
- Negotiates IPsec protocol parameters
- Exchanges public keys
- Authenticates both sides
- Manages keys after the exchange
- Automates entire key-exchange process

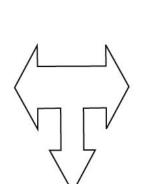


#### Composition of IKE

IKE (Internet Key Exchange) (RFC 2409) Is a Hybrid Protocol

#### SKEME

Mechanism for Utilizing Public Key Encryption for Authentication



#### Oakley

Modes Based Mechanism for Arriving At an Encryption Key Between Two Peers

#### **ISAKMP**

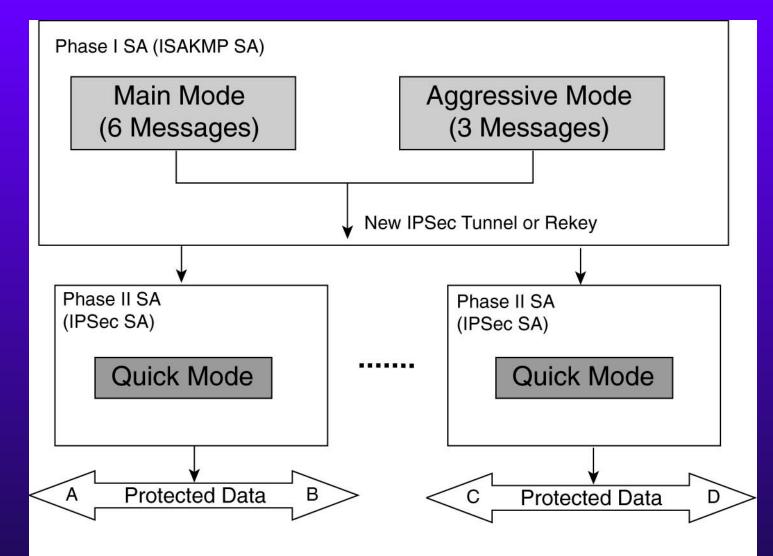
Architecture for Message Exchange Including Packet Formats and State Transitions Between Two Peers

# IPsec Tunnel Creation using IKE

- Identify interesting traffic by an IPsec peer that has been configured to initiate an IPsec session for this traffic
- IPsec peers negotiate a secure authenticated communication channel using main mode or aggressive mode negotiation, resulting in creation of an IKE Security Association (SA) between the two IPsec peers (IKE phase I)
- Create two IPsec SAs between the two IPsec peers via IKE quick mode negotiation (IKE phase II)
- Send data over encrypted tunnel using ESP and/or AH encapsulation



# IKE Main Mode, Aggressive Mode, and Quick Mode



# Goals of Main Mode and Aggressive Mode

- Agreeing on a set of parameters that are to be used to authenticate the two peers
- Agreeing on parameters used to encrypt a portion of the main mode and all of the quick mode messages
- None of the aggressive mode messages are encrypted
- Authenticate the two peers to each other
- Generate keys used to generate keying material for subsequent encryption of data
- All of the parameters negotiated and the keys used to generate keys for encryption are stored as IKE or ISAKMP security association (SA)

## Types of Negotiations by IKE

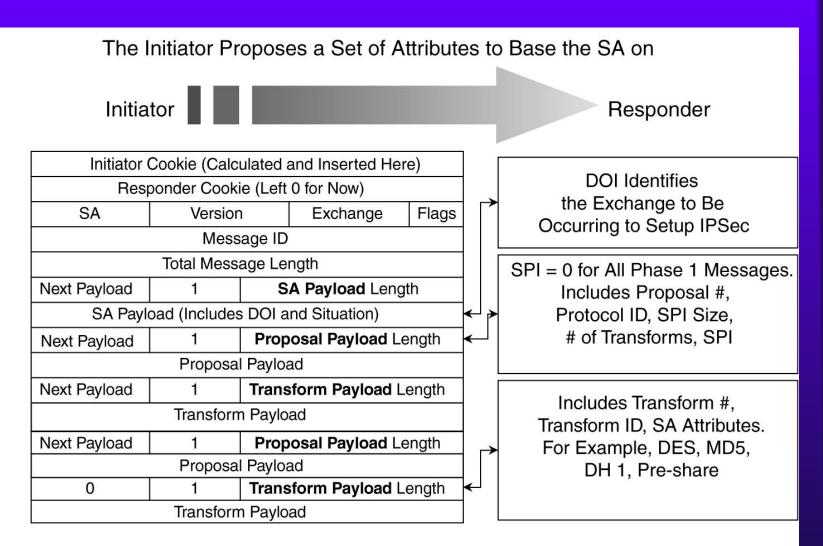
- Main mode using preshared key authentication followed by quick mode negotiation
- Main mode using digital signature authentication followed by quick mode negotiation
- Aggressive mode using preshared key authentication followed by quick mode negotiation
- Main mode using nonces authentication followed by quick mode negotiation
- Aggressive mode using digital signature authentication followed by quick mode negotiation

#### Goals of Quick Mode

 To have two peers agree on a set of attributes for creating the IPsec security associations that could be used by ESP to encrypt the data

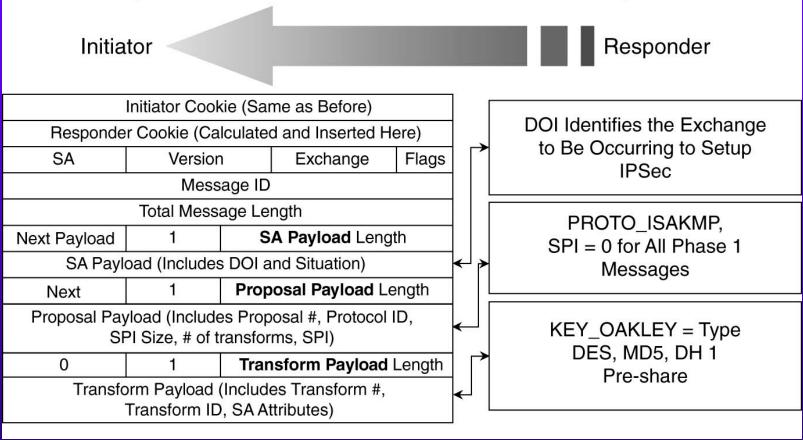
 To redo Diffie-Hellman (DH) exchange so that new keying material can be used to generate IPsec encryption keys

# IKE Main Mode Message 1 using preshared key authentication



# IKE Main Mode Message 2

The Responder Sends Back the One Set of Attributes Acceptable to it



#### Diffie-Hellman Algorithm

 Used in IKE by two peers to generate a shared DH secret and to generate keying material for later use

 DH secret also used with preshared secret to authenticate two peers to each other

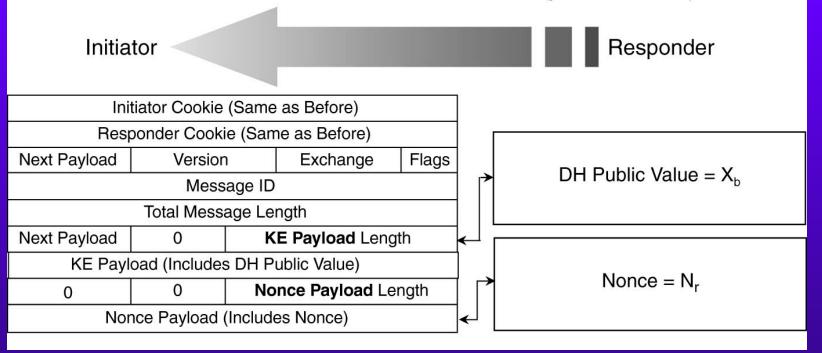
#### Diffie-Hellman Algorithm (cont.)

- There exists X<sub>a</sub> such that X<sub>a</sub> = g<sup>a</sup> mod p where g is the generator, p is a large prime number, and a is a private secret known only to the initiator
- There exists  $X_b$  such that  $X_b = g^b \mod p$  where g is the generator, p is a large prime number, and b is a private secret known only to the responder
- Initiator and responder can generate a shared secret known only to the two of them by exchanging the values X<sub>a</sub> and X<sub>b</sub> with each other
- Initiator secret =  $(X_b)^a \mod p = (X_a)^b \mod p =$ responder secret =  $g^{ab}$

#### IKE Main Mode Message 3 The Initiator Sends Its DH Public Value X<sub>a</sub> and Nonce N<sub>i</sub> Initiator Responder Initiator Cookie (Same as Before) Responder Cookie (Same as Before) Next Payload Version Exchange Flags DH Public Value = $X_a$ Message ID Total Message Length **KE Payload** Length Next Payload 0 ← KE Payload (Includes DH Public Value) Nonce $= N_i$ Nonce Payload Length 0 0 Nonce Payload (Includes Nonce)

## IKE Main Mode Message 4

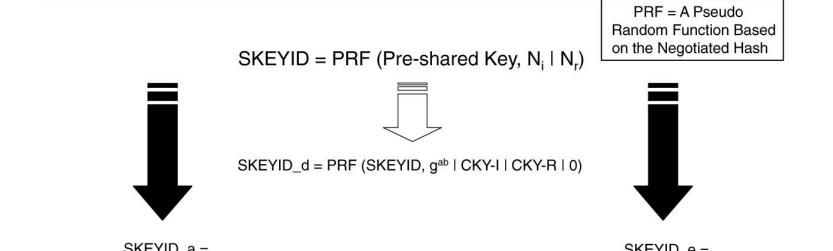
The Resonder Sends Its DH Public Value X<sub>b</sub> and Nonce N<sub>r</sub>



# Session Keys Generated by the Initiator

Calculation of Three Keys (Initiator)

SKEYID\_d — Used to Calculate Subsequent IPSec Keying Material SKEYID\_a — Used to Provide Data Integrity and Authentication to IKE Messages SKEYID\_e — Used to Encrypt IKE Messages



SKEYID\_a = PRF (SKEYID, SKEYID\_d | g<sup>ab</sup> | CKY-I | CKY-R | 1)

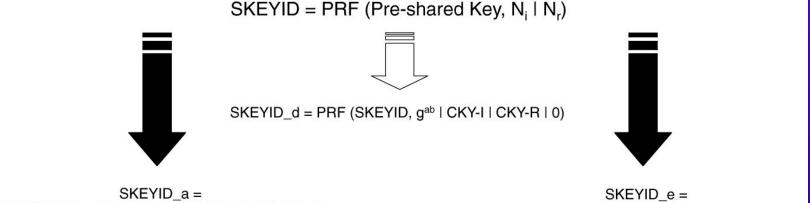
SKEYID\_e = PRF (SKEYID, SKEYID\_a | g<sup>ab</sup> | CKY-I | CKY-R | 2)



# Session Keys Generated by the Responder

Calculation of Three Keys (Responder)

SKEYID\_d — Used to Calculate Subsequent IPSec Keying Material SKEYID\_a — Used to Provide Data Integrity and Authentication to IKE Messages SKEYID\_e — Used to Encrypt IKE Messages

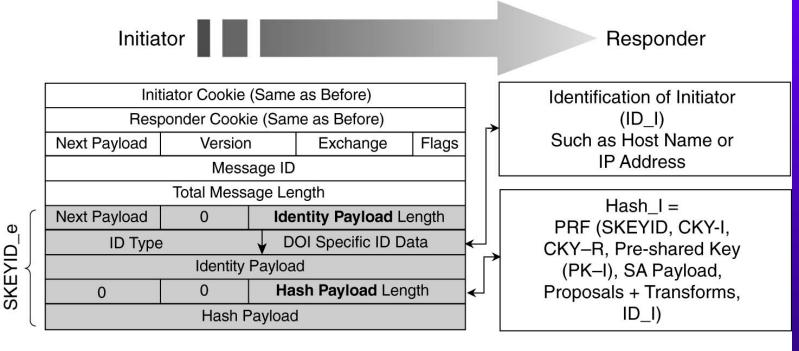


 $SKETID_a =$   $PRF (SKEYID, SKEYID_d | g^{ab} | CKY-I | CKY-R | 1)$ 

SKEYID\_e = PRF (SKEYID, SKEYID\_a | g<sup>ab</sup> | CKY-I | CKY-R | 2)

## IKE Main Mode Message 5

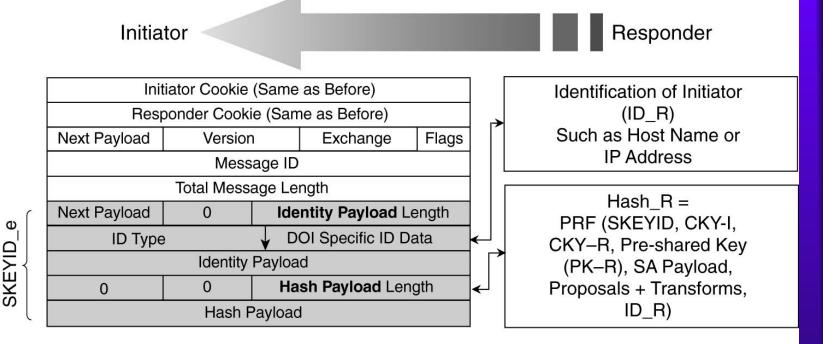
The Initiator Sends Its Authentication Material and ID



-Hash payload and ID\_I are used by responder to authenticate initiator -Identity and hash payloads are encrypted using skeyid\_e

## IKE Main Mode Message 6

The Responder Sends Its Authentication Material and ID



-Hash payload and ID\_R are used by initiator to authenticate responder -Identity and hash payloads are encrypted using skeyid\_e

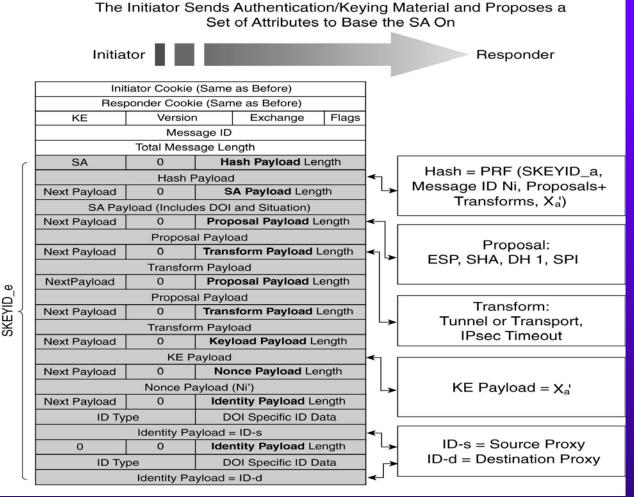
# Completion of IKE Phase I (Main Mode) using Preshared Key

- IKE SA established
- Main mode using preshared key authentication completed
- Quick mode will be used to negotiate parameters of IPsec SA

#### IKE Phase 2 (Quick Mode)

- Negotiate parameters of IPsec SA
- Perfect Forward Secrecy (PFS) may be used by initiator to request that a new DH secret be generated over an encrypted channel
  - New nonces generated:  $N_i$  and  $N_r$
  - New DH public values:
    - X<sub>a</sub>`=g<sup>a</sup> mod p
    - $X_b = g^b \mod p$

### IKE Quick Mode Message



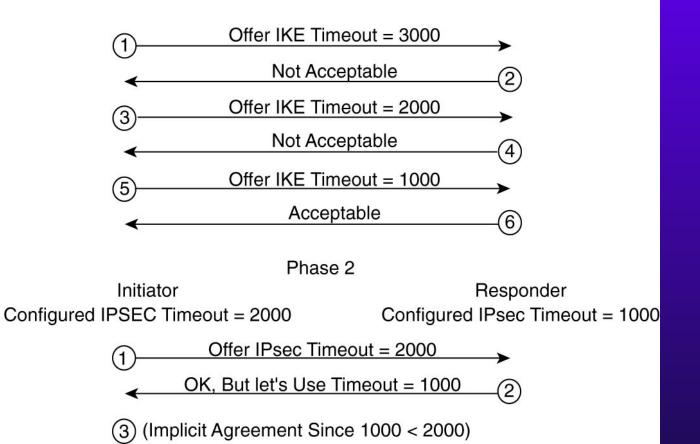
- -Hash used for reauthentication
- -proposal and transform suggests ESP or AH encapsulation type,
- SHA or MD5 integrity checking, DH group, and tunnel or transport mode
- Key exchange payload used for generating new DH secret



# IKE and IPsec Lifetime Negotiation

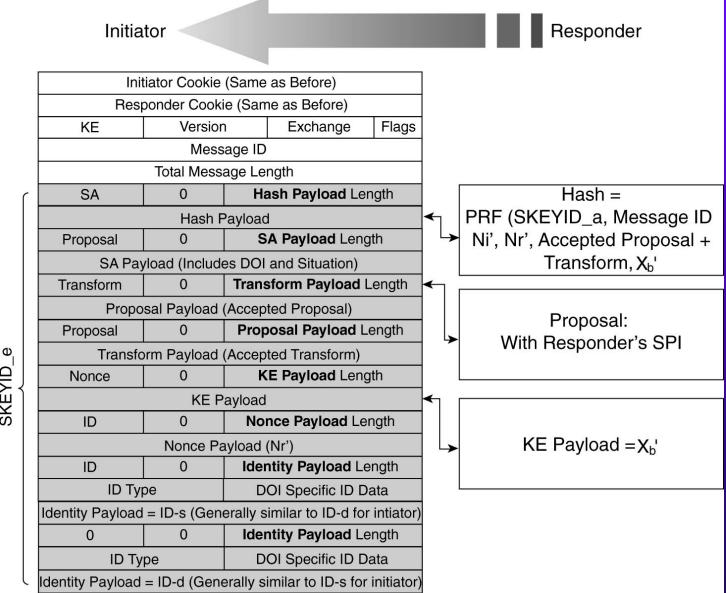
Phase 1

Initiator Responder Configured IKE Timeouts = 3000, 2000, 1000 Configured IKE Timeout = 1500



### IKE Quick Mode Message 2

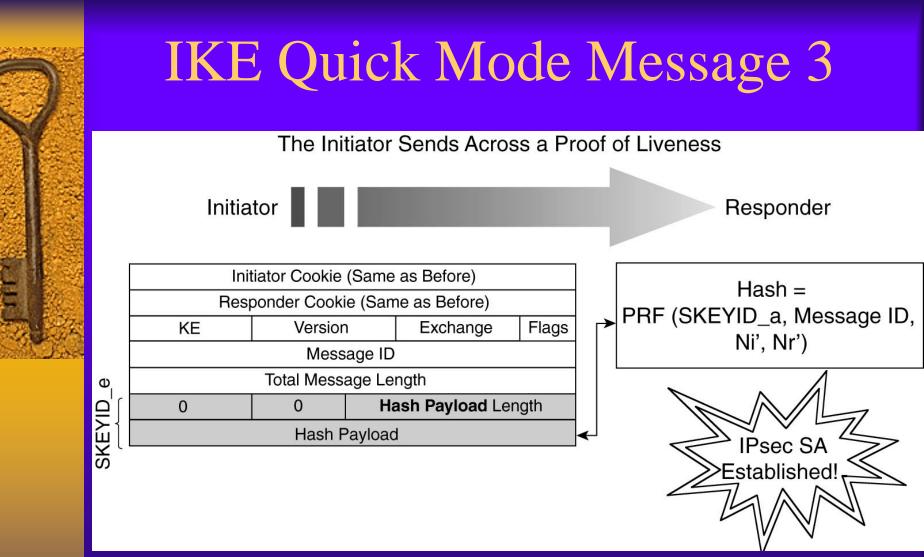
#### The ResponderSends Authentication/Keying Material and Accepts a Set of Attributes to Base the SA On



SKEVID

Generation of IPsec Keying Material

- Both peers generate new DH shared secret =
   (X<sub>b</sub><sup>\*</sup>)<sup>a</sup> mod p = (X<sub>a</sub><sup>\*</sup>)<sup>b</sup> mod p
- Both peers generate shared session keys for incoming and outgoing IPsec SAs based on SKEYID\_d, new DH shared secret, SPI, and N<sub>i</sub>` and N<sub>r</sub>`



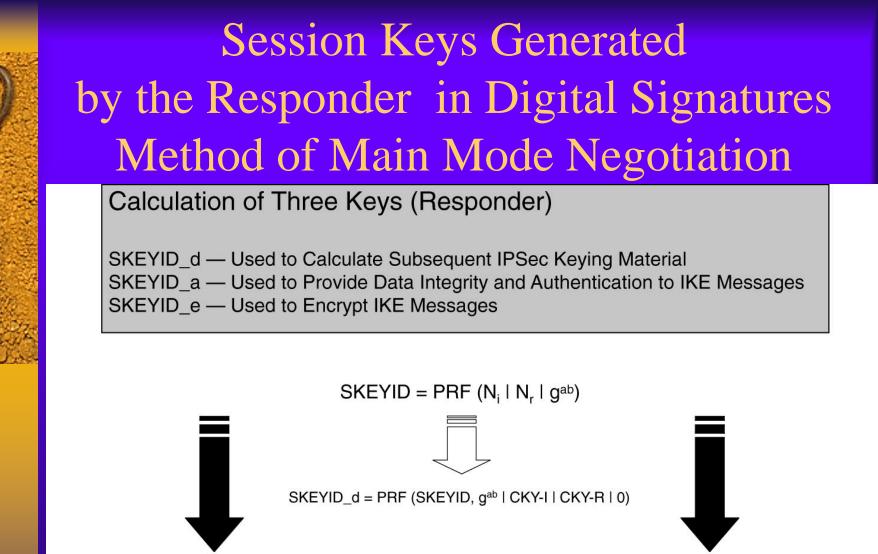
# Main Mode Using Digital Signature Authentication followed by Quick Mode Negotiation

#### Session Keys Generated by the Initiator in Digital Signatures Method of Main Mode Negotiation

Calculation of Three Keys (Initiator)

SKEYID\_d — Used to Calculate Subsequent IPSec Keying Material SKEYID\_a — Used to Provide Data Integrity and Authentication to IKE Messages SKEYID\_e — Used to Encrypt IKE Messages

 $SKEYID_{a} = PRF (SKEYID_{a} = PRF (SKEYID_{a}$ 



SKEYID\_a = PRF (SKEYID, SKEYID\_d | g<sup>ab</sup> | CKY-I | CKY-R | 1)

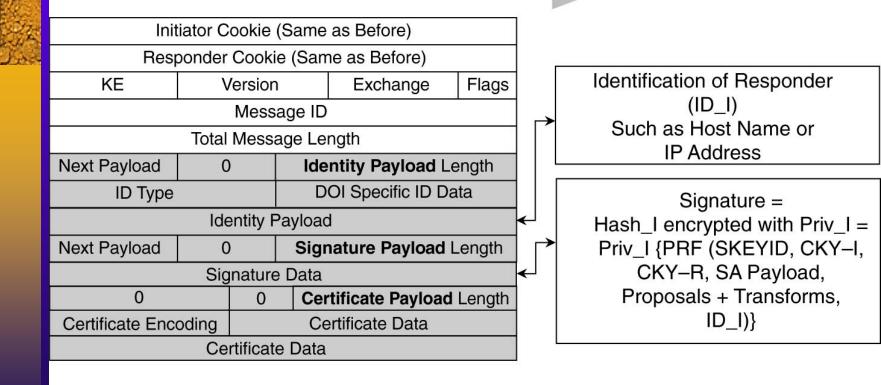
SKEYID\_e = PRF (SKEYID, SKEYID\_a | g<sup>ab</sup> | CKY-I | CKY-R | 2)

# IKE Main Mode Message 5 (using Digital Signatures)

The Initiator Sends Its Authentication Material and ID

Initiator

Responder



# IKE Main Mode Message 6 (using Digital Signatures)

The Responder Sends Its Authentication Material and ID

Initiator

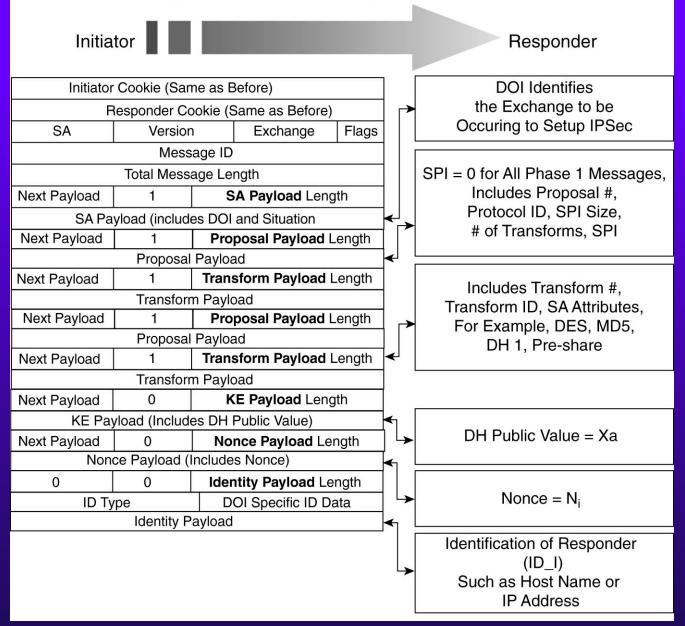
Responder

Initiator Cookie (Same as Before)						]		
Responder Cookie (Same as Before)								
KE	Version			Exchange	Flags	]		Identification of Responder
Message ID							(ID_R)	
Total Message Length							Such as Host Name or IP Address	
Next Payload	0	)	Identity Payload Length				IF Address	
ID Type				DOI Specific ID Data			Signature = Hash_I encrypted with Priv_R =	
Identity Payload						┥		
Next Payload	0		Sig	nature Payload L	ength			Priv_R {PRF (SKEYID, CKY–I,
Signature Data						<b>↓</b>	CKY–R, SA Payload,	
0	0	C	ertificate Payload Length				Proposals + Transforms,	
Certificate Encoding			Certificate Data				ID_R)}	
Certificate Data							L	

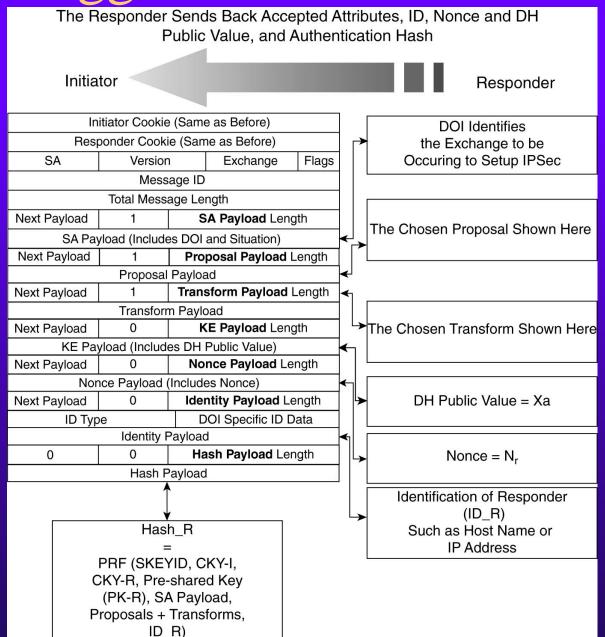
#### Aggressive Mode of IKE Phase 1 using Preshared Key Authentication

### IKE Aggressive Mode Message

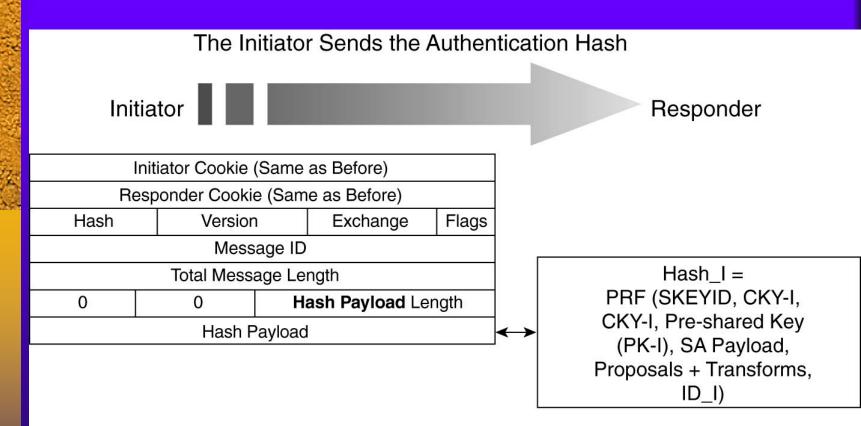
The Initiator Proposes a Set of Attributes, ID, Nonce and DH Public Value



### IKE Aggressive Mode Message 2



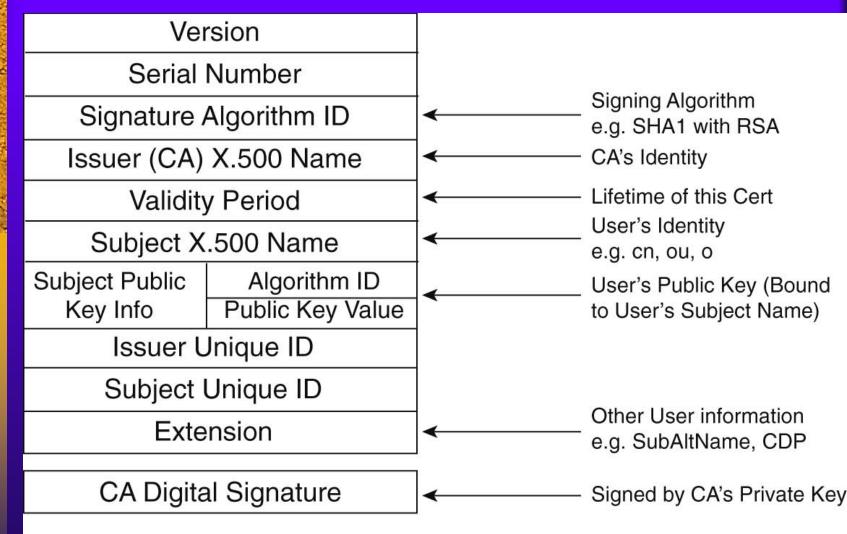
## IKE Aggressive Mode Message 3



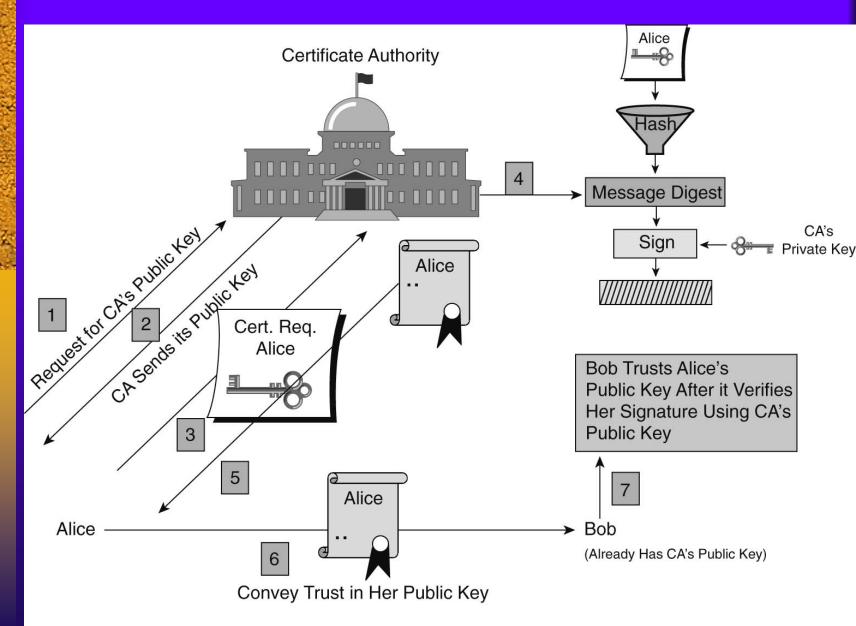
## IKE Device Authentication Methods

- Preshared keys
- Digital signatures
- Encrypted nonces

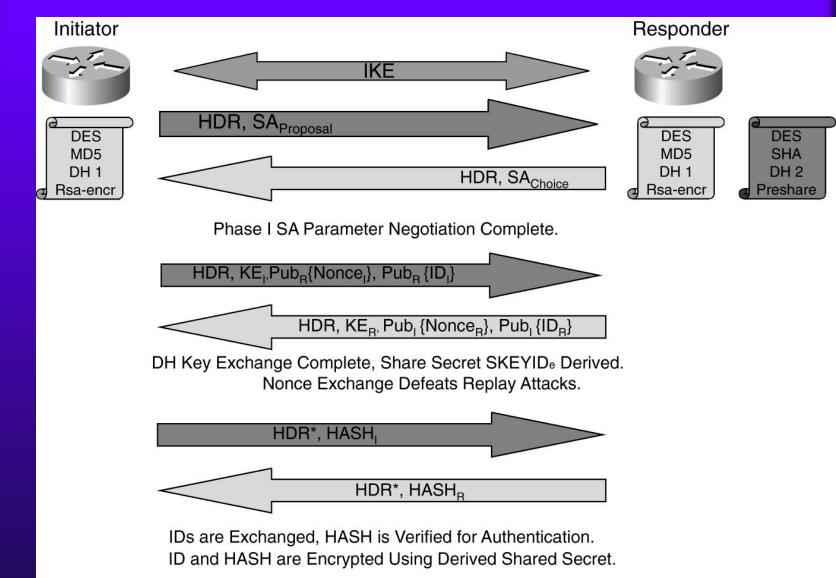
## Contents of a Digital Certificate



## Using Digital Certificates



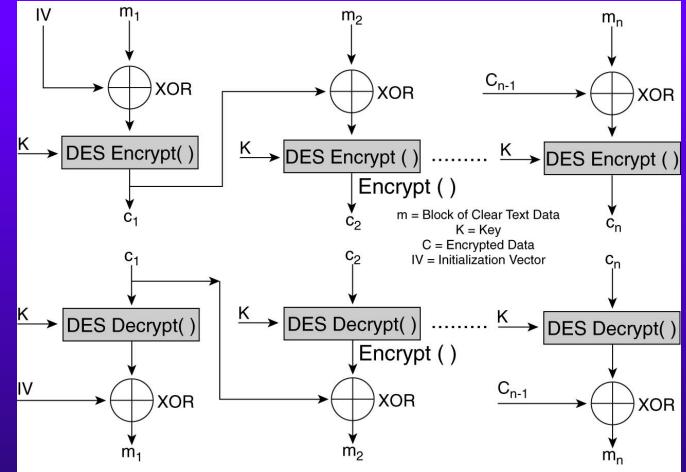
## IKE Main Mode Using Encrypted Nonces



#### **Encryption Methods in IPsec**

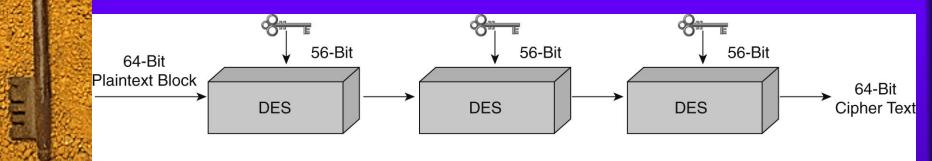
Data Encryption Standard (DES)
Triple DES (3DES)

#### DES Encryption using Cipher Block Chaining (CBC)



Cipher block: DES encryption algorithm converting fixed-length message into cipher text of same length
block size of DES is 64 bits while key length is 56 bits
Initialization vector is sent in ESP header

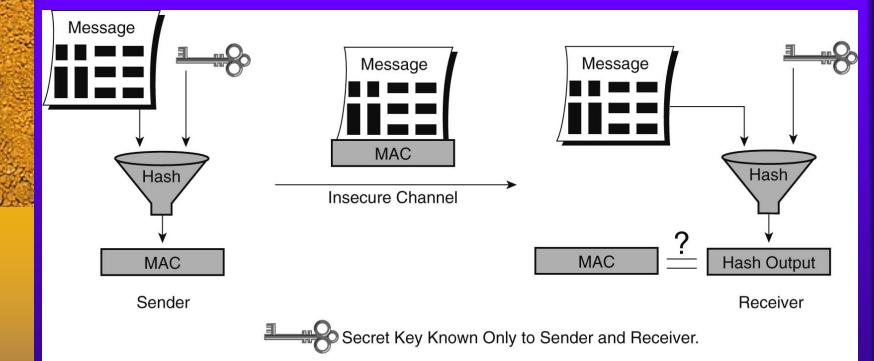
#### **3DES Encryption**



Overall key length is 168 bits

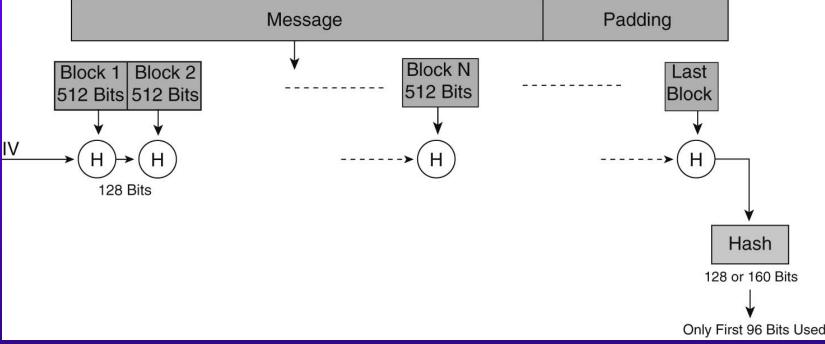
# Integrity Checking Mechanism in IPsec

## Integrity Checking Using Hashes





## Use of Hashes in ESP and AH



MD5 or SHA hashes are truncated to 96 bits



## Packet Encapsulation in IPsec

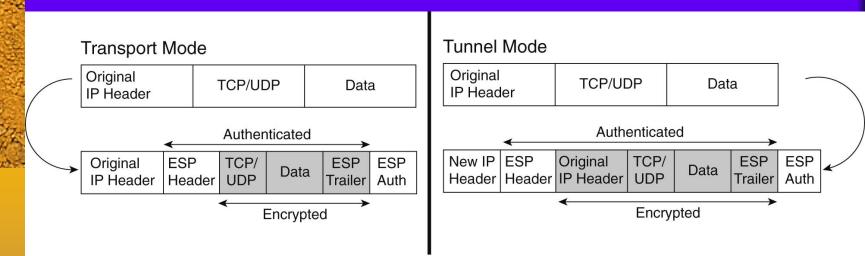
Transport mode

Tunnel mode

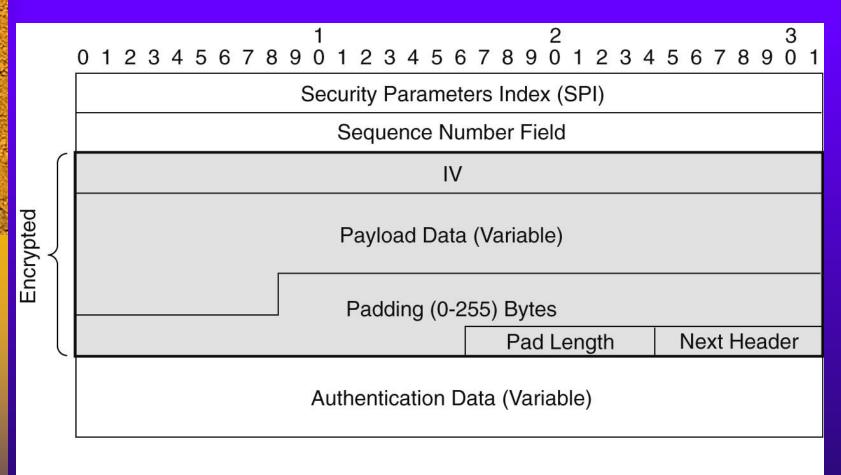
# Packet Format Using AH in Tunnel and Transport Modes

100														
112 X C	Transport N	Node			Tunnel Mode									
	Original IP Header	тс	P/UDP	Data	Original IP Head	Dat	ta							
22.5	Original IP Header	AH	TCP/UD	P Data	New IP Header	AH	Original IP Header	TCP/UDP	Data					
	Aı	henticated E Field in New		<b>→</b>										

# Packet Format Using ESP in Tunnel and Transport Modes



#### **ESP** Header Format



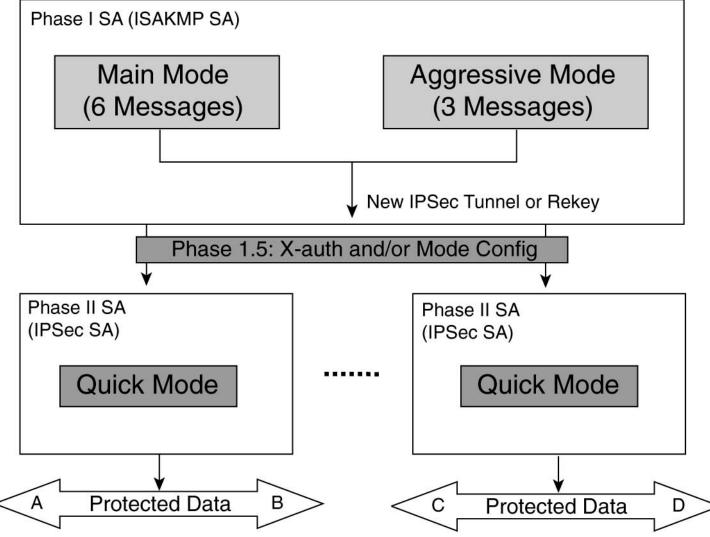
#### AH Header Format

22426	0		2	0	4	-	0	-	0	~	1	4	0	0	4	-	~	7	0	0	2	4	0	~	4	-	0	7	~	0	3	ч
1	0	1 2	2	3	4	5	6	1	8	9	0		2	3	4	5	6	1	8	9	0	1	2	3	4	5	6	1	8	9	0	1
		Next Header Payload Len									RESERVED																					
	Sequence Parameters Index (SPI)																															
	Sequence Number Field																															
8	Authentication Data (Variable)																															

# IKE Enhancements for Remote-Access Client IPsec

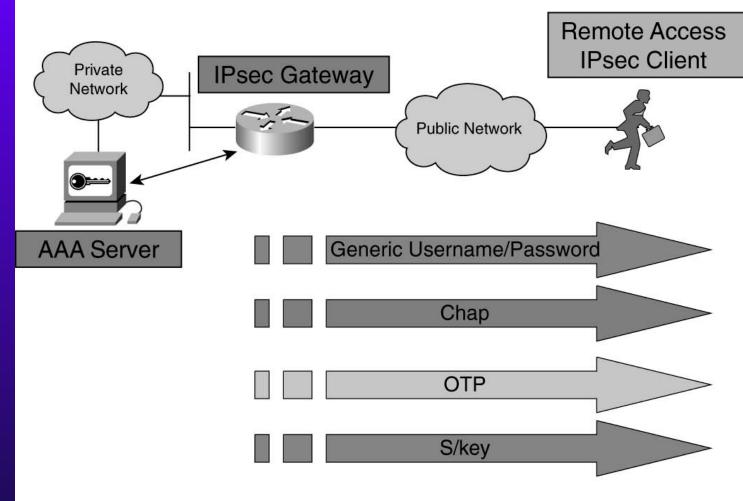


# Extended Authentication and Mode Config in IKE

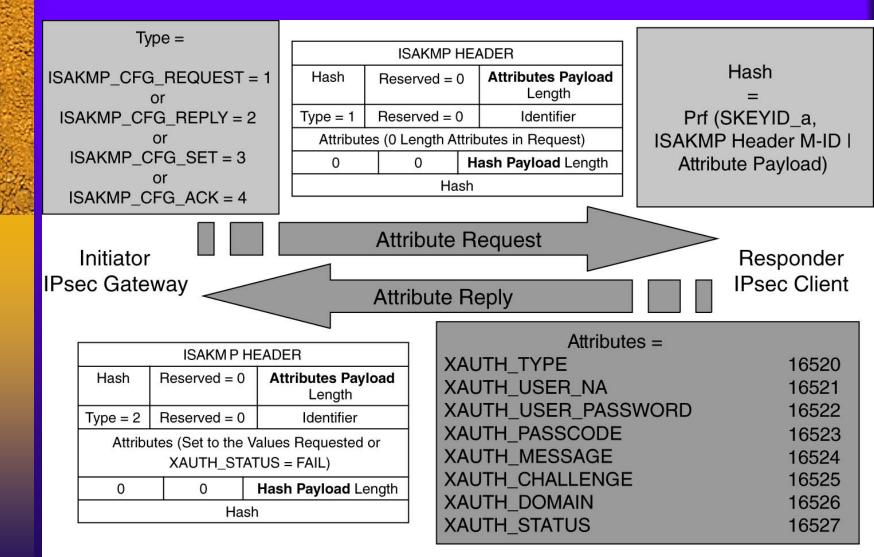


# Negotiation of X-Auth During IKE Negotiation

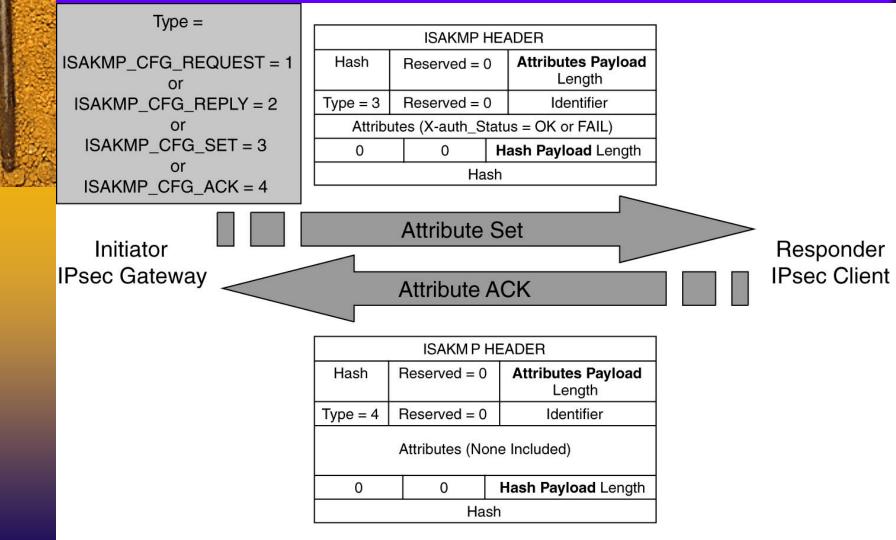
Mechanism Used to Perform Per User Authentication for RA Clients



#### Start of X-Auth with Exchange of Attribute Payloads Using ISAKMP Messages

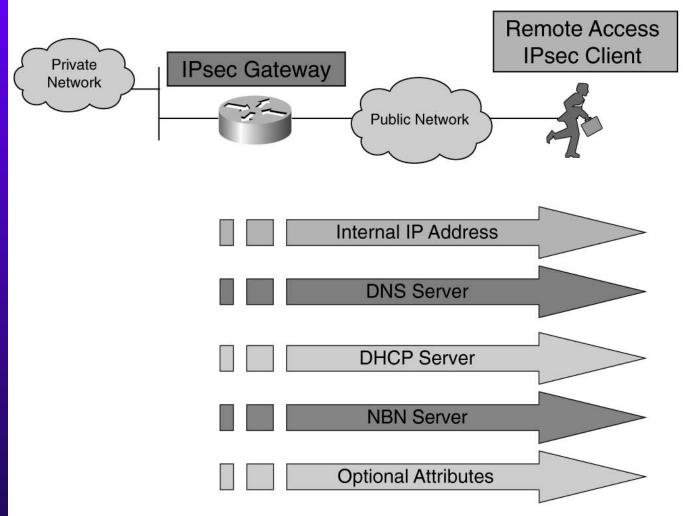


#### Completion of X-Auth with Exchange of Attribute Payloads Using ISAKMP Messages

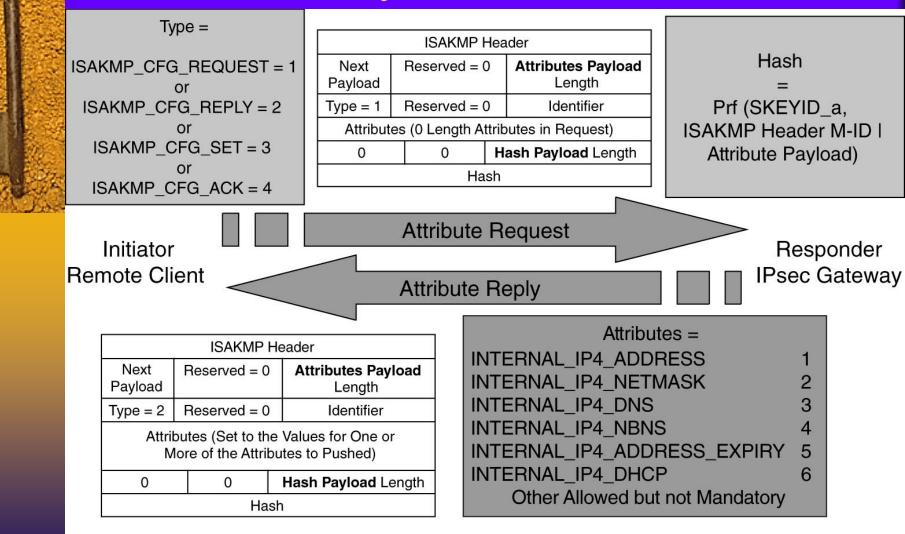


# Mode Configuration During IKE Negotiation

Mechanism Used to Push Attributes to Remote Access IPsec Clients



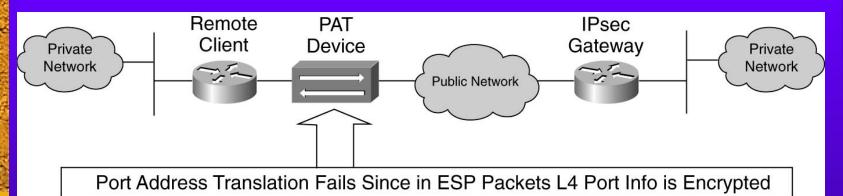
#### Exchanging Attribute Payloads Between the Gateway and the IPsec Client





## NAT Transparency

## Tagging on a UDP Header to Traverse PAT



Ţ		xternal Header	ESP Header	Original IP Header	TCP/UDP Header	Payload	ESP Trailer
External IP Header	24	8 Bytes 0	ESP Header	Original IP Header	TCP/UDP Header	Payload	ESP Trailer